

# A Consideration of an efficient calculation over the extension field of degree 3 and 4 for elliptic curve pairing cryptography

Author

Okayama University

Future University Hakodate

University Rennes 1

Akihiro Sanada

Yuta Kodera

Yasuyuki Nogami

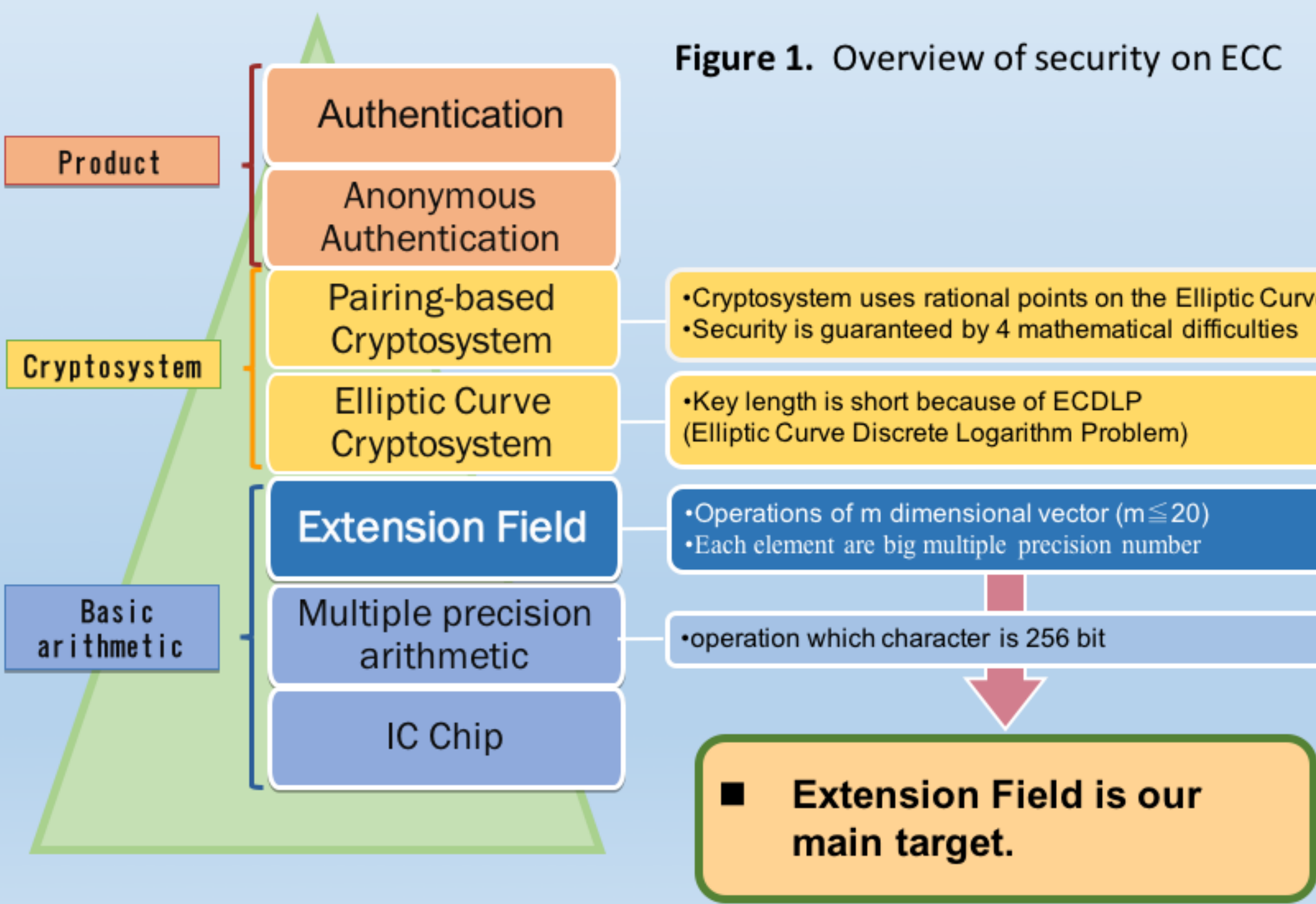
Masaaki Shirase

Sylvain Duquesne

## Introduction

Recently, pairing-based cryptography has been paid much attention since it realizes many innovative security applications. Since recent pairings are defined over extension fields  $F_{p^{18}}$  and  $F_{p^{24}}$ , the efficiency of pairing depends on the arithmetic in these extension field.

Thus, this work considers multiplication, squaring, Frobenius mapping, and inversion over  $F_{p^3}$  and  $F_{p^4}$  as the base of  $F_{p^{18}}$  and  $F_{p^{24}}$  in order to make them more efficient by using CVMA and then evaluate these calculation costs with Karatsuba-based method.



## Example over $F_{p^2}$

### Karatsuba based method

let  $X = a + b\omega, Y = c + d\omega$   
base  $\{1, \omega\} : \omega^2 = -1$

$$\begin{aligned} XY &= (a + b\omega)(c + d\omega) \\ &= ac + (ad + bc)\omega + bd\omega^2 \\ &= ac - db + [(a + b)(c + d) - ac - bd]\omega \end{aligned}$$

### CVMA based (Cyclic Vector Multiplication Algorithm)

let  $X = a\omega + b\omega^2, Y = c\omega + d\omega^2$   
base  $\{\omega, \omega^2\} : \omega^3 = 1 = -\omega - \omega^2$

$$\begin{aligned} XY &= (a\omega + b\omega^2)(c\omega + d\omega^2) \\ &= ac\omega^2 + (ad + bc)\omega^3 + bd\omega^4 \\ &= \{(a - b)(c - d) - ac\}\omega + \{(a - b)(c - d) - bd\}\omega^2 \end{aligned}$$

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## Result of degree 3

As a competitive construction, Karatsuba-based method uses polynomial basis  $\{1, \alpha, \alpha^2\}$  defined by method uses polynomial  $f(x) = x^3 - 2$  and its zero  $\alpha \in F_{p^3}$ . In this case, it needs to satisfy  $2^{\frac{p-1}{3}} \not\equiv 1 \pmod{p}$  so that  $f(x)$  becomes irreducible over  $F_p$ .

On the other hand, the proposed method uses the basis  $\{\tau_1, \tau_2, \tau_3\} = \{\omega - \omega^{-1}, \omega^2 - \omega^{-2}, \omega^3 - \omega^{-3}\}$  defined by modular polynomial  $\Phi_7(x) = (x^7 - 1)/(x - 1)$  and its zero  $\omega \in F_{p^3}$ . In this case,  $p \not\equiv 1, 6 \pmod{p}$  such that  $\{\tau_1, \tau_2, \tau_3\}$  becomes a basis. It is noted that our method is available for more prime numbers as  $p$  than conventional method, moreover it realizes efficient multiplication, squaring, and Frobenius mapping as shown in Figure 2.

### Karatsuba based method's multiplication over $F_{p^3}$

$$\begin{aligned} &(a_0 + a_1\alpha + a_2\alpha^2)(b_0 + b_1\alpha + b_2\alpha^2) \\ &= a_0b_0 + 2[(a_1 + a_2)(b_1 + b_2) - a_1b_1 - a_2b_2] \\ &\quad + [(a_0 + a_1)(b_0 + b_1) - a_0b_0 - a_1b_1 + 2a_2b_2]\alpha \\ &\quad + [(a_0 + a_2)(b_0 + b_2) - a_0b_0 - a_2b_2 + a_1b_1]\alpha^2 \end{aligned}$$

$$1: a_0b_0 + 2[(a_1 + a_2)(b_1 + b_2) - a_1b_1 - a_2b_2]$$

$$\alpha: (a_0 + a_1)(b_0 + b_1) - a_0b_0 - a_1b_1 + 2a_2b_2$$

$$\alpha^2: (a_0 + a_2)(b_0 + b_2) - a_0b_0 - a_2b_2 + a_1b_1$$

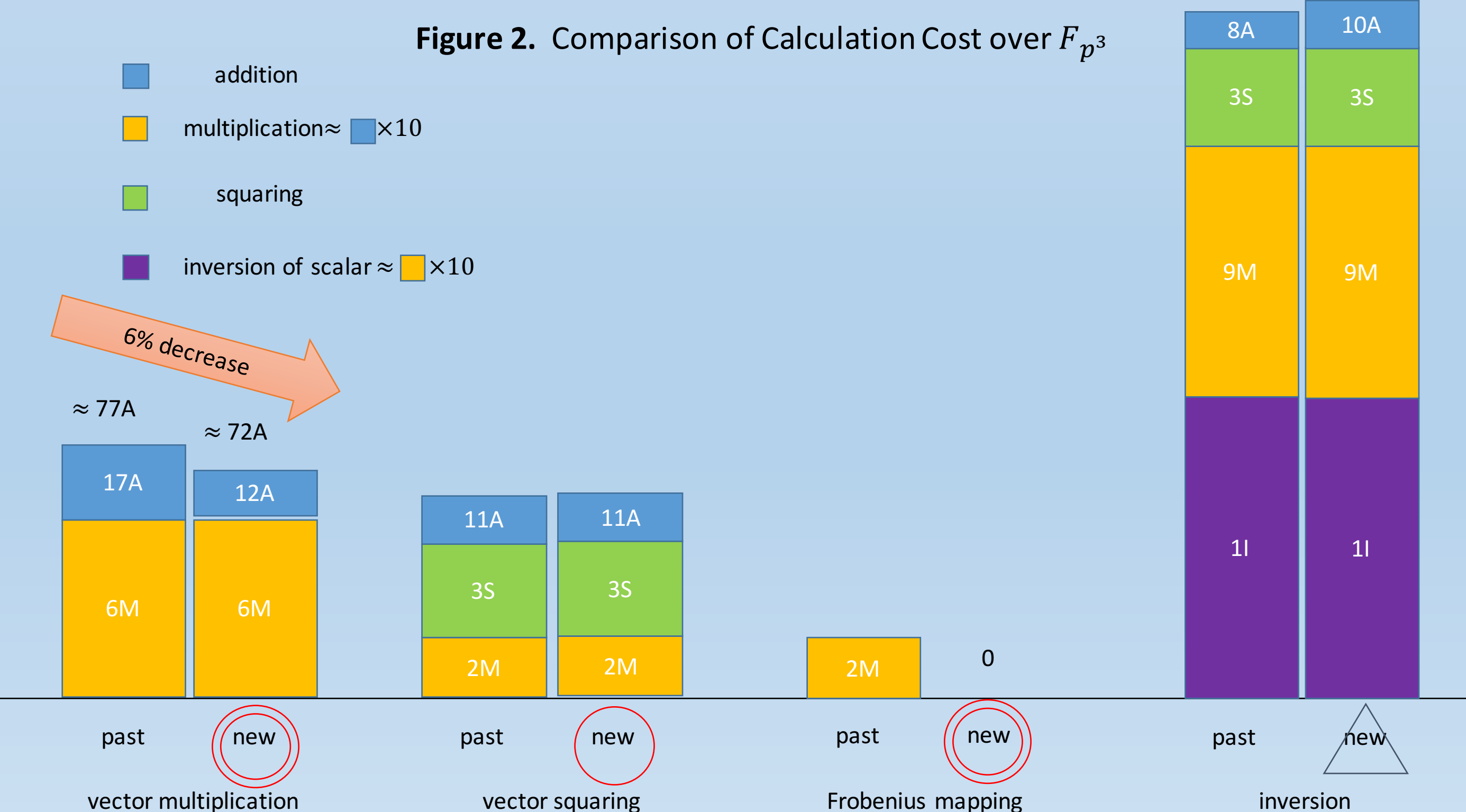
### CVMA based method's multiplication over $F_{p^3}$

$$(a_1\tau_1 + a_2\tau_2 + a_3\tau_3)(b_1\tau_1 + b_2\tau_2 + b_3\tau_3)$$

$$\tau_1: (a_1 - a_2)(b_2 - b_1) + (a_2 - a_3)(b_3 - b_2) - a_1b_1$$

$$\tau_2: (a_2 - a_3)(b_3 - b_2) + (a_1 - a_3)(b_3 - b_1) - a_2b_2$$

$$\tau_3: (a_1 - a_2)(b_2 - b_1) + (a_2 - a_3)(b_3 - b_2) - a_1b_1$$



## Result of degree 4

Efficient extension field  $F_{p^4}$  with Karatsuba-based method is constructed by towering technique such as  $F_{(p^2)^2}$ . Modular polynomial are  $f(x) = x^2 + 1$  and  $g(x) = x^2 - (\omega + 1)$ , where  $\omega$  is zero of  $f(x)$ . Then,  $F_{(p^2)^2}$  is constructed by the basis  $\{1, \omega\} \times \{1, \theta\}$ , where  $\theta$  is a zero of  $g(x)$ . For this construction,  $p$  needs to satisfy  $p \equiv 3 \pmod{8}$ .

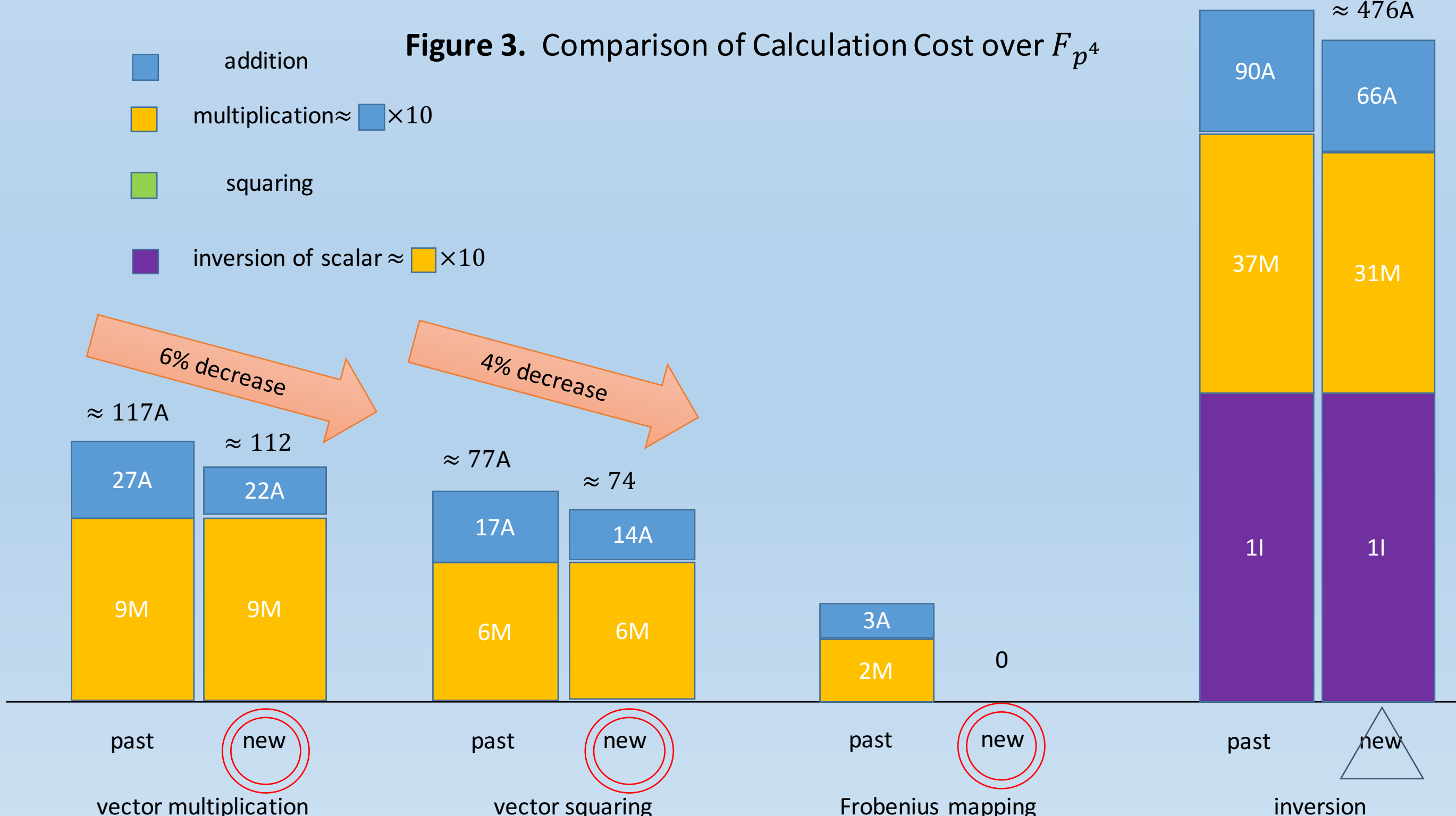
The proposed method uses the basis  $\{\omega, \omega^2, \omega^3, \omega^4\}$  defined by modular polynomial  $\Phi_5(x) = (x^5 - 1)/(x - 1)$  and its zero  $\omega \in F_{p^4}$ . In this case,  $p \equiv 2, 3 \pmod{5}$  such that  $\Phi_5(x)$  becomes irreducible over  $F_p$ . The proposed CVMA-based method provides an efficient calculation as shown in Figure 3.

### CVMA based method's multiplication over $F_{p^4}$

$$\begin{aligned} &(x_0\omega + x_1\omega^2 + x_2\omega^3 + x_3\omega^4)(y_0\omega + y_1\omega^2 + y_2\omega^3 + y_3\omega^4) \\ &= (x_2y_2 + x_1y_3 - x_3y_1 - x_0y_3 - x_1y_2 + x_2y_1 - x_3y_0)\omega \\ &\quad + (x_0y_0 + x_2y_3 - x_3y_2 - x_0y_3 - x_1y_2 + x_2y_1 - x_3y_0)\omega^2 \\ &\quad + (x_3y_3 + x_0y_1 - x_1y_0 - x_0y_3 - x_1y_2 + x_2y_1 - x_3y_0)\omega^3 \\ &\quad + (x_1y_1 + x_0y_2 - x_2y_0 - x_0y_3 - x_1y_2 + x_2y_1 - x_3y_0)\omega^4 \\ &= \{U - (x_1 - x_3)(y_1 - y_3) - x_0y_0\}\omega \\ &\quad + \{U - (x_2 - x_3)(y_2 - y_3) - x_1y_1\}\omega^2 \\ &\quad + \{U - (x_0 - x_1)(y_0 - y_1) - x_2y_2\}\omega^3 \\ &\quad + \{U - (x_0 - x_2)(y_0 - y_2) - x_3y_3\}\omega^4 \end{aligned}$$

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$$\begin{aligned} &= (x_0 - x_3)(y_0 - y_3) + (x_1 - x_2)(y_1 - y_2) \\ &= (x_0 + x_1 - x_2 - x_3)(y_0 + y_1 - y_2 - y_3)\{(x_0 - x_3)(y_1 - y_2) + (x_1 - x_2)(y_0 - y_3)\} \\ &= (x_0 + x_1 - x_2 - x_3)(y_0 + y_1 - y_2 - y_3) \\ &\quad + (x_0 - x_1)(y_0 - y_1) - (x_0 - x_2)(y_0 - y_2) \\ &\quad - (x_1 - x_3)(y_1 - y_3) + (x_2 - x_3)(y_2 - y_3) \end{aligned}$$



## Conclusions

Most of CVMA-based arithmetic are superior to those of the Karatsuba-based method. It is noted that the modular prime number  $p$  for each condition needs to satisfy a certain condition. When  $p$  satisfies both of them, our proposed method has better costs.

**Table 1.** Comparison of calculation costs

Operation	Karatsuba-based	CVMA-based
Comparison over $F_{p^3}$		
Multiplication	6M+17A	6M+12A
Squaring	2M+3S+11A	2M+3S+11A
Frobenius mapping	2M	0
Inversion	9M+3S+8A+I	9M+3S+10A+I
Comparison over $F_{p^4}$		
Multiplication	9M+27A	9M+22A
Squaring	6M+17A	6M+14A
Frobenius mapping	2M+3A	0
Inversion	37M+90A+I	31M+66A+I

## Literature cited

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## Future work

We will apply the proposed methods in pairing based Cryptosystem (ex. Figure 4) to evaluate the efficiency of the proposed methods.

**Figure 4.** Cryptosystem based on Math

